

FM-Indexing Grammars Induced by Suffix Sorting for Long Patterns

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Outline

- Text Index
- FM-Index
- Geometric-BWT
- Our Index
- Experimental Result
- Future Work

Text Index

- Given a text T over alphabet Σ , a text index is a data structure built from T supporting the following queries with pattern P .
- $\text{exist}(P)$: Does P occur in T ?
- $\text{count}(P)$: How often does P occur in T ?
- $\text{locate}(P)$: Where does P occur in T ?

0123456789012
 $T = \text{bananabana\$}$
 $P = \text{ana}$

Occurrences:
bananabana\$
bananabana\$
bananabana\$
bananabana\$

$\text{exist}(P) = \text{true}$
 $\text{count}(P) = 4$
 $\text{locate}(P) = (1, 3, 7, 9)$

$T = \text{bananabana\$}$

FM-Index

[Ferragina & Manzini JACM'05]

↓
 n rotations
bananabana\$
ananabana\$b
nanabana\$ba
anabana\$ban
nabana\$bana
abana\$banan
banana\$banana
anana\$bananab
nana\$bananaba
ana\$bananaban
na\$bananabana
a\$bananabana
\$bananabana

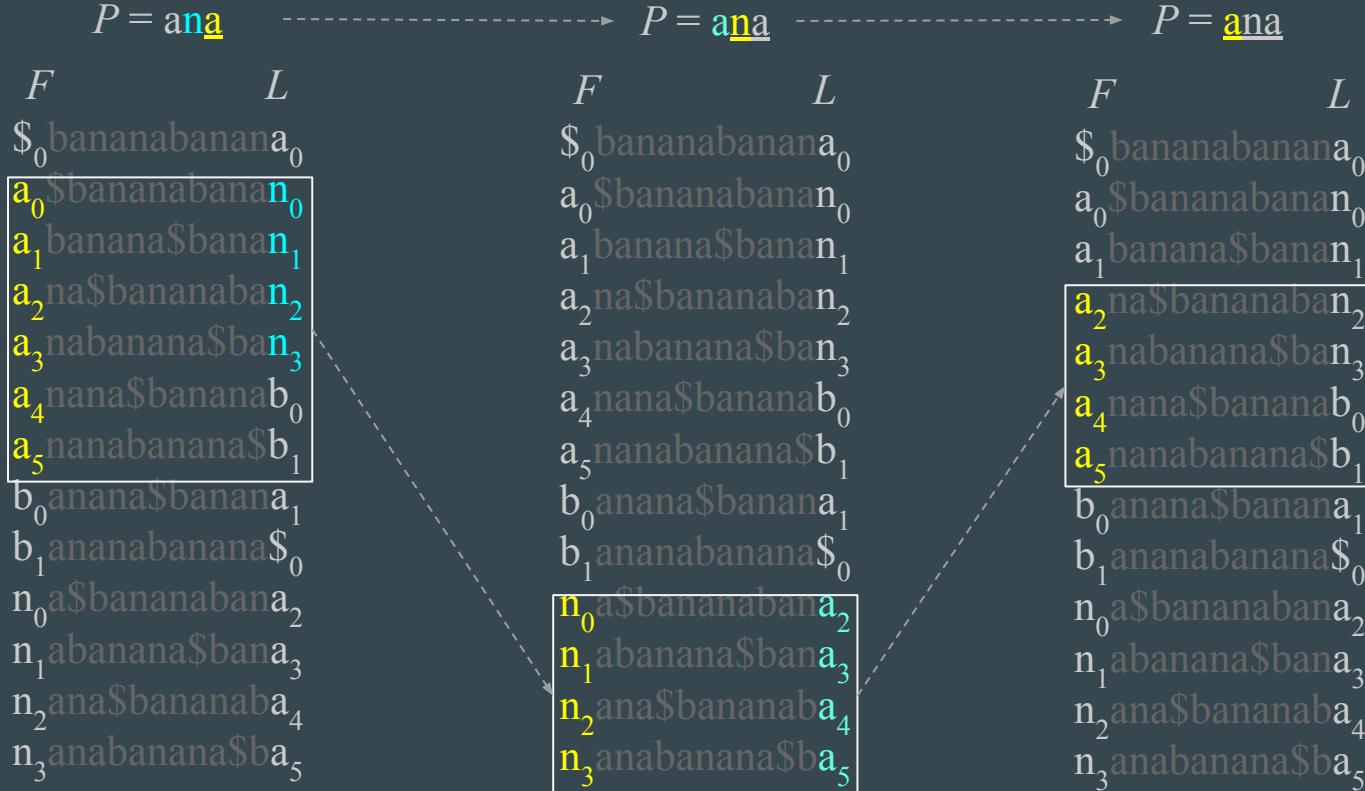
sort

\$bananabana
a\$bananabana
abanana\$banan
ana\$bananaban
anabana\$ban
anana\$bananab
ananabana\$b
banana\$banana
bananabana\$
na\$bananabana
nabana\$bana
nana\$bananaba
nanabana\$ba

store only
information of
first (F) & last (L)
columns

F	L
\$bananabana	a\$bananabana
a\$bananabana	abanana\$banan
abanana\$banan	ana\$bananaban
ana\$bananaban	ana\$bananaban
anabana\$ban	anabana\$ban
anana\$bananab	anana\$bananab
ananabana\$b	ananabana\$b
banana\$banana	banana\$banana
bananabana\$	bananabana\$
na\$bananabana	na\$bananabana
nabana\$bana	nabana\$bana
nana\$bananaba	nana\$bananaba
nanabana\$ba	nanabana\$ba

FM-Index (Counting by Backward-Searching)



FM-Index (LF-Mapping)

- $\text{begin}(F, c)$: where is begin of maximal-character-run of character c ?
- $\text{rank}(L, i, c)$: How many character c is in prefix $L[0 : i-1]$ of string L ?
- LF-Mapping(i): $\text{begin}(F, c) + \text{rank}(L, i, c)$
- Counting P on FM-index is reduced to $O(|P|)$ LF-Mappings.
- $\text{begin}(F, c)$ can be supported by simple prefix-sum array.
- $\text{rank}(L, i, c)$ can be support by Wavelet Tree [Grossi et al, SODA'03]
- Disadvantage: LF-Mapping is not cache-friendly.

Geometric BWT

[Yu-Feng Cieng et al. Algorithmica'05]

- λ -factorization: Viewing λ consecutive characters as single meta-character (fitted into single machine word).
- $T^{(\lambda)}$: T follows λ -factorization
- $P^{(\lambda, k)}$: P follows λ -factorization and the last factor is of length k , $1 \leq k \leq \lambda$
- Reducing pattern counting to backward searching + 4-sided range counting.

$$\begin{array}{ccc} T = \text{bananabananan}\$ & \xrightarrow{\lambda = 2} & T^{(2)} = \text{ba.na.na.ba.na.na.\$\$} \\ P = \text{anana} & & P^{(2, 1)} = \text{an.an.a?} \\ & & P^{(2, 2)} = ?\text{a.na.na} \end{array}$$

Geometric BWT

$T^{(2)} = \text{ba.na.na.ba.na.na}.\$\$$

colexgraphical rank range

$P^{(2,1)} = \text{an.an.a?}$

$F^{(2)}$	$L^{(2)}$
$\$\$_0$	na_0
ba_0	na_1
ba_1	$\$\$_0$
na_0	na_2
na_1	na_3
na_2	ba_0
na_3	ba_1

There is no a? in $F^{(2)}$

$P^{(2,2)} = ?\text{a.na.na}$

$F^{(2)}$	$L^{(2)}$
$\$\$_0$	na_0
ba_0	na_1
ba_1	$\$\$_0$
na_0	na_2
na_1	na_3
na_2	ba_0
na_3	ba_1

$P^{(2,2)} = ?\underline{\text{a}}.\underline{\text{na}}.\underline{\text{na}}$

$F^{(2)}$	$L^{(2)}$
$\$\$_0$	na_0
ba_0	na_1
ba_1	$\$\$_0$
na_0	na_2
na_1	na_3
na_2	ba_0
na_3	ba_1

$P^{(2,2)} = ?\underline{\text{a}}.\underline{\text{na}}.\underline{\text{na}}$

$F^{(2)}$	$L^{(2)}$
$\$\$_0$	na_0
ba_0	na_1
ba_1	$\$\$_0$
na_0	na_2
na_1	na_3
na_2	ba_0
na_3	ba_1

suffix (array offset) range of na.na

Geometric-BWT

- $O(|P|)$ LF-Mapping on FM-index = $O(|P|/\lambda)$ LF-Mapping + 4-sided range counting on Geometric BWT λ times.
- Disadvantage: There are at most λ factorizations of pattern matching genuine occurrences on text.

Our Index (LMS-Factorization)

- LMS is short for LeftMost S-type (S^*), which is a term from one of the linear-time suffix array construction algorithms, SAIS [Ge Nong et al, DCC'09], and later an grammar compression scheme, GCIS [D.S.N Nunes et al, DCC'18].

SL-types = L L S* S S S L L L L S*

$T^{(LMS)}$ = b | a n | a n | a b | a n | a n a | \$
SL-types = L | S* L L | S*



Our Index (LMS-Factorization)

- There are at most 2 different factorizations of P , $P^{(\text{LMS}, S)}$ and $P^{(\text{LMS}, L)}$.

$$P = \text{anana}$$

$$\begin{array}{ll} P_S^{(\text{LMS})} &= \text{a n} \mid \text{a n} \mid \text{a} \\ \text{SL-types} &= \text{L L} \mid \text{S}^* \text{L} \mid \text{S}^* \end{array}$$

$$\begin{array}{ll} P_L^{(\text{LMS})} &= \text{a n} \mid \text{a n a} \\ \text{SL-types} &= \text{L L} \mid \text{S}^* \text{L L} \end{array}$$

Our Index (LMS-Factorization + λ -factorization)

- LMS-factors could be long substring of T , we further apply λ -factorization on each LMS-factor of T .
- The lex(icographical) & colex(icographical) grammar rules, G_{lex} and G_{colex} , are mapping factors in set of factors to their lex. and colex. rank.

$T = \text{bananabana} \text{\$}$

$\lambda = 2$

G_{lex}	G_{colex}
$0 : \$$	$0 : \$$
$1 : a$	$1 : a$
$2 : ab$	$2 : b$
$3 : an$	$3 : ba$
$4 : b$	$4 : na$

$T^{(\text{LMS}, 2)} = b \mid a \mid n \mid a \mid n \mid [a \mid b] \mid a \mid n \mid [a \mid n] \cdot a \mid \$$
SL-types = $L \mid S^* \mid L \cdot L \mid S^*$

$T^{(\text{LMS}, 2, G\text{lex})} = 4 \ 3 \ 3 \ 2 \ 3 \ 3 \ 1$

θ

For simplicity, we take $T^{(\text{LMS}, 2)}$ for the purpose of presentation

Our Index (LMS-Factorization + λ -factorization)

- There are at most 2 different LMS-factorizations of P , $P_S^{(LMS)}$ and $P_L^{(LMS)}$.
- There are at most λ factorizations of first LMS-factor of $P_S^{(LMS)}$ and $P_L^{(LMS)}$.

$$\begin{array}{c} P = \text{anana} \\ \swarrow \quad \searrow \\ \begin{array}{l} P_S^{(LMS, 1)} = a . n \mid a n \mid a \\ P_S^{(LMS, 2)} = a n \mid a n \mid a \\ \text{SL-types} = L \ L \mid S^* L \mid S^* \end{array} \\ \begin{array}{l} P_L^{(LMS, 1)} = a . n \mid a n . a \\ P_L^{(LMS, 2)} = a n \mid a n . a \\ \text{SL-types} = L \ L \mid S^* L \ L \end{array} \end{array}$$

Note that $P_S^{(LMS, 1)} = P_L^{(LMS, 1)}$ and $P_S^{(LMS, 2)} = P_L^{(LMS, 2)}$ in terms of final factorization in this sample pattern, but it is not the case in general.

Our Index (Counting Multiple-LMS-factor Pattern)

$T = \text{bananabana}\$$

$P = \text{anana}$

$\lambda = 2$

$T^{(\text{LMS}, 2)} = b | a n | a n | a b | a n | a n . a | \$$

$$\begin{array}{c} P_S^{(\text{LMS}, 1)} = * a . n | \text{a n} | \underline{\text{a}} \\ \xrightarrow{* -} F^{(\text{LMS}, 2)} L^{(\text{LMS}, 2)} \end{array} \quad \begin{array}{c} P_S^{(\text{LMS}, 1)} = * a . \text{n} | \underline{\text{a n}} | \underline{a *} \\ \xrightarrow{F^{(\text{LMS}, 2)} L^{(\text{LMS}, 2)}} \end{array}$$

$\$_0$	a_0
a_0	an_0
ab_0	an_1
an_0	an_2
an_1	an_3
an_2	ab_0
an_3	b_0
b_0	$\$_0$

$\$_0$	a_0
a_0	an_0
ab_0	an_1
an_0	an_2
an_1	an_3
an_2	ab_0
an_3	b_0
b_0	$\$_0$

There is no n in the range

Our Index (Counting Multiple-LMS-factor Pattern)

$T = \text{bananabana}\$$

$P = \text{anana}$

$\lambda = 2$

$T^{(\text{LMS}, 2)} = b | a n | a n | a b | a n | a n . a | \$$

$P_S^{(\text{LMS}, 2)} = * a n | \text{a n} | \underline{\text{a}*} \longrightarrow P_S^{(\text{LMS}, 2)} = * \text{a n} | \underline{\text{a n}} | \underline{\text{a}*}$

$F^{(\text{LMS}, 2)} L^{(\text{LMS}, 2)}$

$\$_0$	a_0
a_0	an_0
ab_0	an_1
an_0	an_2
an_1	an_3
an_2	ab_0
an_3	b_0
b_0	$\$_0$

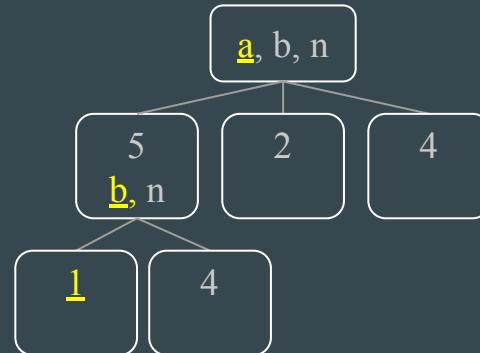
$F^{(\text{LMS}, 2)} L^{(\text{LMS}, 2)}$

$\$_0$	a_0
a_0	an_0
ab_0	an_1
an_0	an_2
an_1	an_3
an_2	ab_0
an_3	b_0
b_0	$\$_0$

Our Index (Counting Single-LMS-factor Pattern)

- If LMS-factorization of pattern results in only single LMS-factor and λ -factorization on the LMS-factor results in multiple factors, counting is following the algorithm applied on Geometric-BWT.
- If final factorization of pattern is single factor, there is a generalized suffix trie, GST , built from all factors on $T^{(LMS, \lambda)}$ along with count information in nodes of the GST to deal with such scenario.

$T^{(LMS, 2)} = b | a n | a n | a b | a n | a n . a |$
\$
 $P = ab$



Experimental Result (Dataset)

- Cere, E.Coli, and Para from [Repetitive Corpus on Pizza & Chili Corpus](#)
- Chr19.15 is concatenation of 15 different human chromosome 19 processed from [chr19.1000.fa on tudocomp dataset](#)
- Artificial.x are artificial texts generated for simulating concatenated generational genome sequences, where the mutation rate is x%.

Experimental Result (Construction Time)

dataset	FM-index	RLFM ⁽⁰⁾		RLFM ⁽¹⁾	
		time [s]	time [s]	λ	time [s]
CERE	101.6	102.6	1	881.3	
			2	471.4	
			3	365.0	
			4	299.4	
			5	290.4	
			6	277.4	
			7	275.9	
			8	276.6	
CHR19.15	207.7	208.9	1	2093.7	
			2	1107.7	
			3	807.3	
			4	683.0	
			5	633.2	
			6	626.6	
			7	609.6	
			8	597.2	
E.COLI	24.5	25.0	1	187.1	
			2	110.3	
			3	87.1	
			4	71.9	
			5	66.7	
			6	65.8	
			7	66.3	
			8	71.1	
PARA	98.0	98.0	1	755.7	
			2	410.5	
			3	319.3	
			4	267.4	
			5	260.9	
			6	251.5	
			7	248.2	
			8	252.2	

dataset	FM-index	RLFM ⁽⁰⁾		RLFM ⁽¹⁾	
		time [s]	time [s]	λ	time [s]
ARTIFICIAL.1	130.0	131.4	1	616.1	
			2	339.4	
			3	263.2	
			4	227.4	
			5	224.1	
			6	230.7	
			7	230.5	
			8	233.7	
ARTIFICIAL.2	129.2	132.9	1	578.2	
			2	328.3	
			3	254.5	
			4	218.7	
			5	213.2	
			6	220.5	
			7	220.3	
			8	224.2	
ARTIFICIAL.4	130.5	137.1	1	555.2	
			2	555.2	
			3	243.4	
			4	213.6	
			5	209.7	
			6	218.1	
			7	216.5	
			8	222.8	
ARTIFICIAL.8	132.4	144.4	1	530.6	
			2	306.4	
			3	241.0	
			4	211.0	
			5	207.1	
			6	214.9	
			7	215.5	
			8	219.7	

RLFM [Mäkinen & Navarro CPM'05]

Experimental Result (Space)

input text			RLFM ⁽⁰⁾				RLFM ⁽¹⁾			
name	space [MiB]	σ	$r^{(0)}$ [M]	space [MiB]	λ	space [MiB]	$\sigma^{(1)}$	$r^{(1)}$ [M]	lg P	
CERE	439.9	6	11.6	26.8	1	26.5	6	11.6	-	
					2	20.4	26	8.3	-	
					3	18.4	95	6.7	12	
					4	17.3	271	5.8	11	
					5	16.7	602	5.3	11	
					6	15.1	1081	5.1	12	
					7	14.9	1790	5.0	13	
					8	14.9	2810	4.9	13	
CHR19.15	845.8	6	32.3	70.8	1	69.7	6	32.3	-	
					2	54.9	22	23.5	-	
					3	50.8	58	19.0	10	
					4	47.1	140	16.5	9	
					5	44.9	305	15.1	-	
					6	40.3	620	14.3	11	
					7	39.7	1174	13.8	12	
					8	39.4	2086	13.6	13	
E.COLI	107.5	16	15.0	26.2	1	25.4	16	15.0	-	
					2	20.4	123	10.7	-	
					3	19.3	399	8.5	-	
					4	17.8	809	7.3	13	
					5	17.2	1272	6.7	12	
					6	15.3	1764	6.3	13	
					7	15.1	2356	6.2	13	
					8	15.1	3251	6.1	-	
PARA	409.4	6	15.6	34.4	1	34.0	6	15.6	-	
					2	26.5	26	11.3	-	
					3	24.5	96	9.1	12	
					4	22.6	296	7.9	11	
					5	20.0	774	7.2	11	
					6	19.6	1620	6.9	12	
					7	19.4	2701	6.7	13	
					8	19.3	4013	6.7	14	

input text			RLFM ⁽⁰⁾				RLFM ⁽¹⁾			
name	space [MiB]	σ	$r^{(0)}$ [M]	space [MiB]	λ	space [MiB]	$\sigma^{(1)}$	$r^{(1)}$ [M]	lg P	
ARTIFICIAL.1	502.5	5	50.9	91.4	1	89.5	5	50.9	-	
					2	84.9	17	39.0	8	
					3	71.2	51	32.3	7	
					4	69.2	131	28.7	8	
					5	68.0	297	26.7	9	
					6	67.4	611	25.8	10	
					7	67.1	1164	25.3	11	
					8	66.9	2058	25.1	11	
ARTIFICIAL.2	500.0	5	87.5	141.8	1	138.5	5	87.5	-	
					2	131.5	17	67.0	7	
					3	111.8	51	55.5	7	
					4	109.9	131	49.3	7	
					5	108.5	297	45.9	9	
					6	107.8	611	44.3	9	
					7	107.4	1164	43.5	10	
					8	107.2	2069	43.2	11	
ARTIFICIAL.4	495.0	5	147.0	215.8	1	210.1	5	147.0	-	
					2	198.7	17	111.4	6	
					3	169.7	51	91.8	6	
					4	168.4	131	81.1	7	
					5	167.0	297	75.4	8	
					6	166.0	611	72.6	9	
					7	165.4	1164	71.3	10	
					8	165.0	2077	70.7	11	
ARTIFICIAL.8	485.0	5	237.4	300.2	1	290.9	5	237.4	-	
					2	286.4	17	174.3	8	
					3	239.6	51	141.3	7	
					4	235.1	131	123.4	7	
					5	231.0	297	113.9	8	
					6	228.4	611	109.2	9	
					7	226.6	1164	107.0	10	
					8	225.5	2079	106.1	11	

Future Work

- Is λ -factorization required?
- How to augment the current index to support $\text{locate}(P)$?
- Comparison of $\text{count}(P)$ to FM-index family (e.g. RLFM+ [Sirén et al. SPIRE'08], Faster-Minuter [Gog et al. DCC'16])
- Comparison of $\text{locate}(P)$ to GCIIS family (e.g. Akagi et al, SPIRE'21, Díaz et al. SPIRE'21)

Feedback or Question?